

Reversal Ratio and Linear Extension Diameter

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- 1 Linear Extension Diameter
- 2 A Constant Bound?
- 3 Posets of fixed dimension
- 4 Posets of fixed width

Definition

Let \mathbf{P} be a finite poset. The *linear extension graph* $G(\mathbf{P}) = (V, E)$ of \mathbf{P} is defined as follows:

- V is the set of all linear extensions of \mathbf{P} and
- two linear extensions are adjacent if and only if they differ only in the transposition of a single (adjacent) pair of points.

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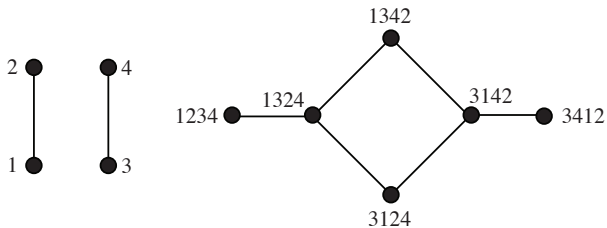
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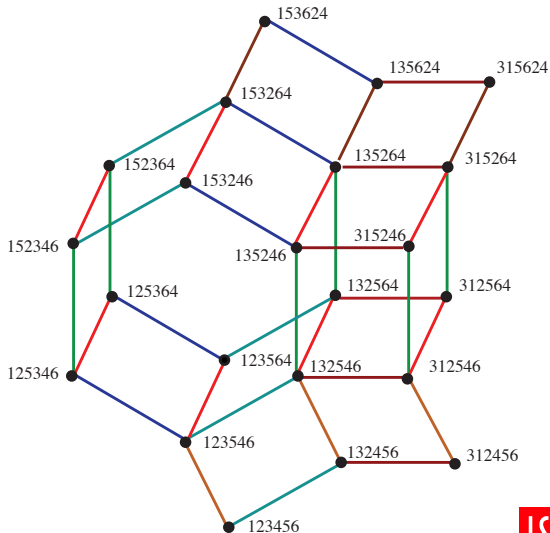
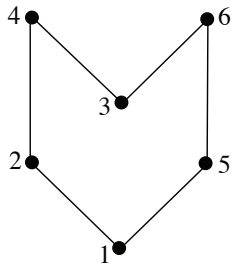
Definition (Felsner and Reuter 1999)

The *linear extension diameter* of a finite poset \mathbf{P} , denoted $\text{led}(\mathbf{P})$, is the diameter of its linear extension graph $G(\mathbf{P})$.

Example



Another Example



Felsner and Massow (2011)



Definition

Let \mathbf{P} be a poset and L_1, L_2 linear extensions of \mathbf{P} . We define the *reversal ratio of the pair* (L_1, L_2) as

$$RR(\mathbf{P}; L_1, L_2) = \frac{\text{dist}(L_1, L_2)}{\text{inc}(\mathbf{P})}.$$

The *reversal ratio of* \mathbf{P} is

$$RR(\mathbf{P}) = \frac{\text{led}(\mathbf{P})}{\text{inc}(\mathbf{P})} = \max_{L_1, L_2} RR(\mathbf{P}; L_1, L_2).$$

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 - Unpublished example difficult to analyze.



How small can $RR(\mathbf{P})$ be?

Theorem (BK 2011+)

For every sufficiently large positive integer k , there exists a poset \mathbf{P}_k of width k with $RR(\mathbf{P}_k) \leq C/\log k$.

Doubling Property

Definition

Let $\mathbf{G} = (A \cup B, E)$ be a bipartite graph with $|A| = |B| = k$. We say that \mathbf{G} has the doubling property if for every $Y \subset A$ with $|Y| \leq k/3$,
 $|N(A)| \geq 2|A|$.



Doubling Property

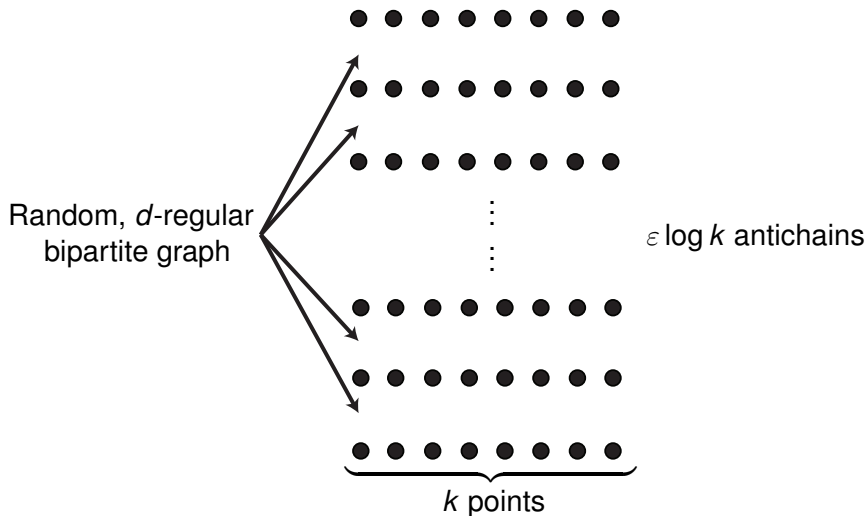
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Lemma

Let $G_d(A, B)$ be a random d -regular bipartite graph on vertex sets A and B of size k , chosen according to the configuration model. For each $d \geq 10$ and k sufficiently large, $G_d(A, B)$ has the doubling property with high probability.

Construction of \mathbf{P}_k



Proposition

The number of incomparable pairs in \mathbf{P}_k is at least

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For k sufficiently large,

$$\text{led}(\mathbf{P}_k) \leq \frac{19}{3}\varepsilon k^2 \log k.$$

Definition

For $d \geq 2$, define $g_{RR}(d) = \inf\{RR(\mathbf{P}) : \dim(\mathbf{P}) = d\}$.

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- Standard example \mathbf{S}_n ? $RR(\mathbf{S}_n) \rightarrow 1$
- d -dimensional grid? $RR(\mathbf{n}^d) \rightarrow 1/2$ (Felsner-Massow)
- $g_{RR}(3) = 2/3$ by considering \mathbf{n}^3
- $1/2 \leq g_{RR}(4) \leq 4/7$

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Corollary

There is a constant C such that $g_{RR}(d) \leq C/\log d$ for all $d \geq 2$.

Definition

For $w \geq 2$, define $h_{RR}(w) = \inf\{RR(\mathbf{P}) : \text{width}(\mathbf{P}) = w\}$.

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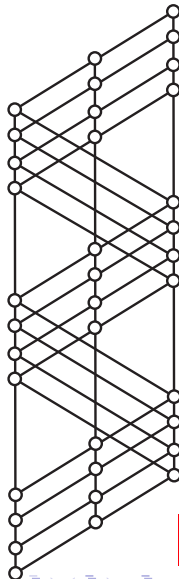
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Proposition

$$h_{RR}(3) \leq 5/6$$

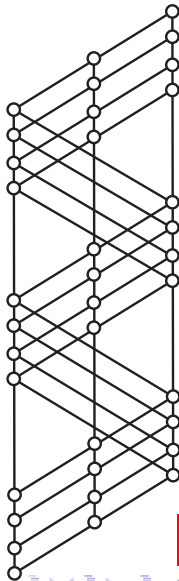
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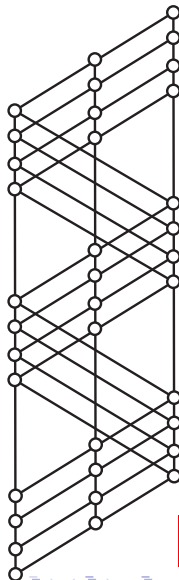
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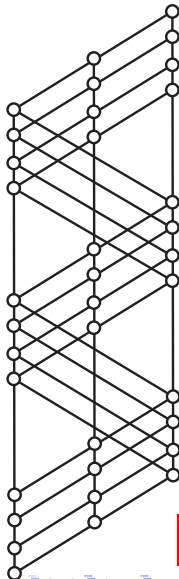


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L_1	L_2
C_4	B_4/C_4
$C_3/A_4/B_4$	A_4
B_3	$A_3/B_3/C_3$
$B_2/C_2/A_3$	C_2
A_2	$C_1/A_2/B_2$
$A_1/B_1/C_1$	B_1
	A_1



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Conjecture (BK)

$$h_{RR}(3) = 3/4$$

Thank You

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